

CIS 771: Software Specifications

Lecture: Alloy Logic (part E)

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Outline

- Declarations
- Set multiplicities
- Relational multiplicities
- Declaration constraints
- Nested multiplicities
- Cardinality constraints

Declarations

A *declaration* introduces a name for a relation whose value is a subset of the value of the *bounding expression* appearing after the ":"

`relation-name : expression`

Examples

Note: the bounding expression is usually formed with unary relations and the arrow operator, but any expression can be used.

- `address : Name -> Addr`
 - maps names to addresses (representing a single address book)
- `addr: Book -> Name -> Addr`
 - maps books to names to addresses (representing a collection of address books)
- `address: Name -> (Name + Addr)`
 - maps names to names and addresses (representing a multilevel address book)

Declarations

A *declaration* introduces a name for a relation whose value is a subset of the value of the *bounding expression* appearing after the ":"

```
relation-name : expression
```

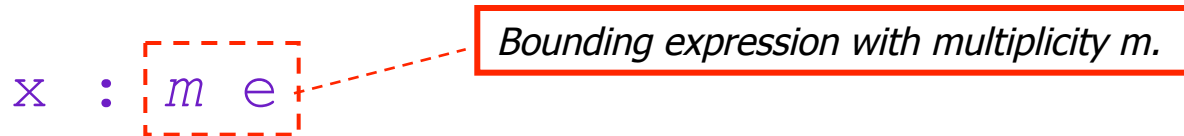
Examples -- with more complicated expressions

- `address : (Alias + Group) -> (Addr + Alias + Group)`
 - maps aliases and groups to addresses, aliases, and groups
- `address : (Alias->Group) + (Group->(Addr + Alias + Group))`
 - has the same "type" as the declaration above
 - aliases, groups in domain; addresses, aliases, groups in co-domain
 - more precise than the version above because it constrains aliases to only map to groups
 - illustrates how declaration expressions can combine relations (e.g., via the union operator) as well as sets

Set Multiplicities

If the *bounding expression* of a declaration denotes a set (is unary), it can be prefixed by a multiplicity keyword m which constrains the size of set x according to m .

$x : m e$



Bounding expression with multiplicity m .

Multiplicity Keywords -- similar to those used for quantification

- **set** any number
- **one** exactly one
- **lone** zero or one
- **some** one or more

For set-value bounded expressions, omitting the multiplicity keyword is the same as writing **one**.

Set Multiplicities

If the *bounding expression* of a declaration denotes a set (is unary), it can be prefixed by a multiplicity keyword m which constrains the size of set x according to m .

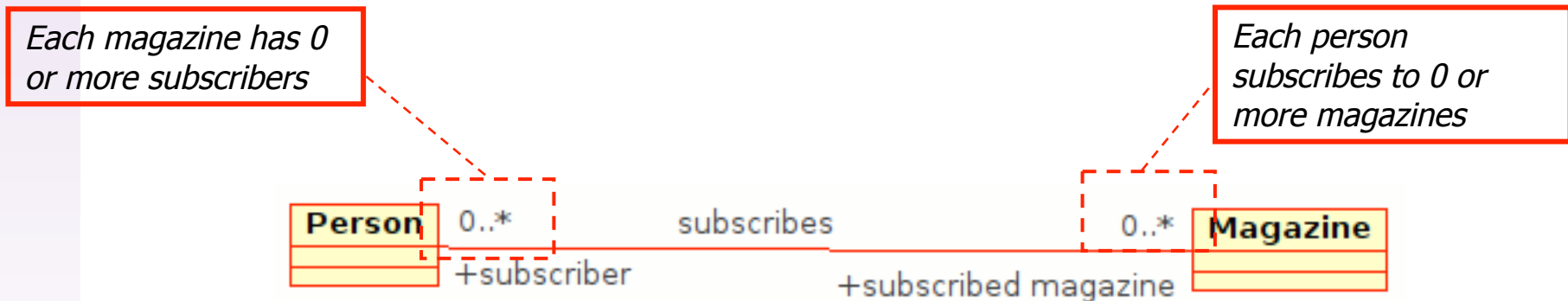
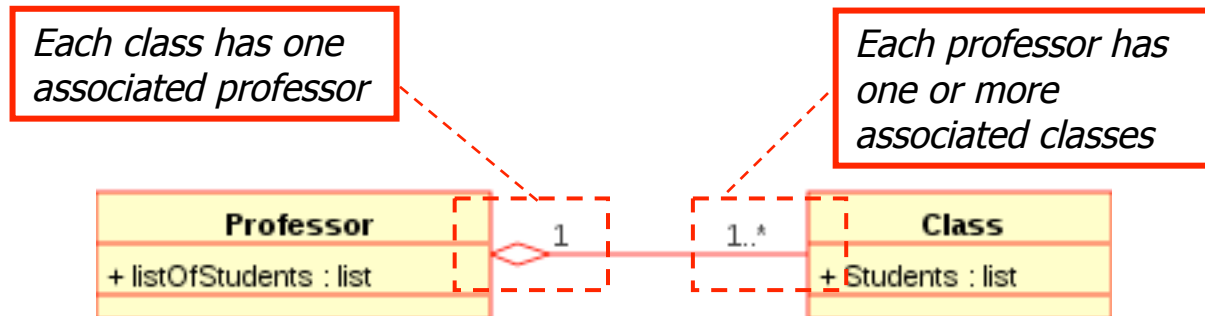
$x : m e$

Examples

- `RecentlyUsed: set Name`
 - says that `RecentlyUsed` is a subset of the set `Name`
- `senderAddress: Addr`
 - says that `senderAddress` is a scaler in the set `Addr`
- `senderName: lone Name`
 - says that `senderName` is an option: either a scaler in the set `Name`, or empty
- `receiverAddresses: some Addr`
 - says that `receiverAddresses` is a non-empty subset of `Addr`.

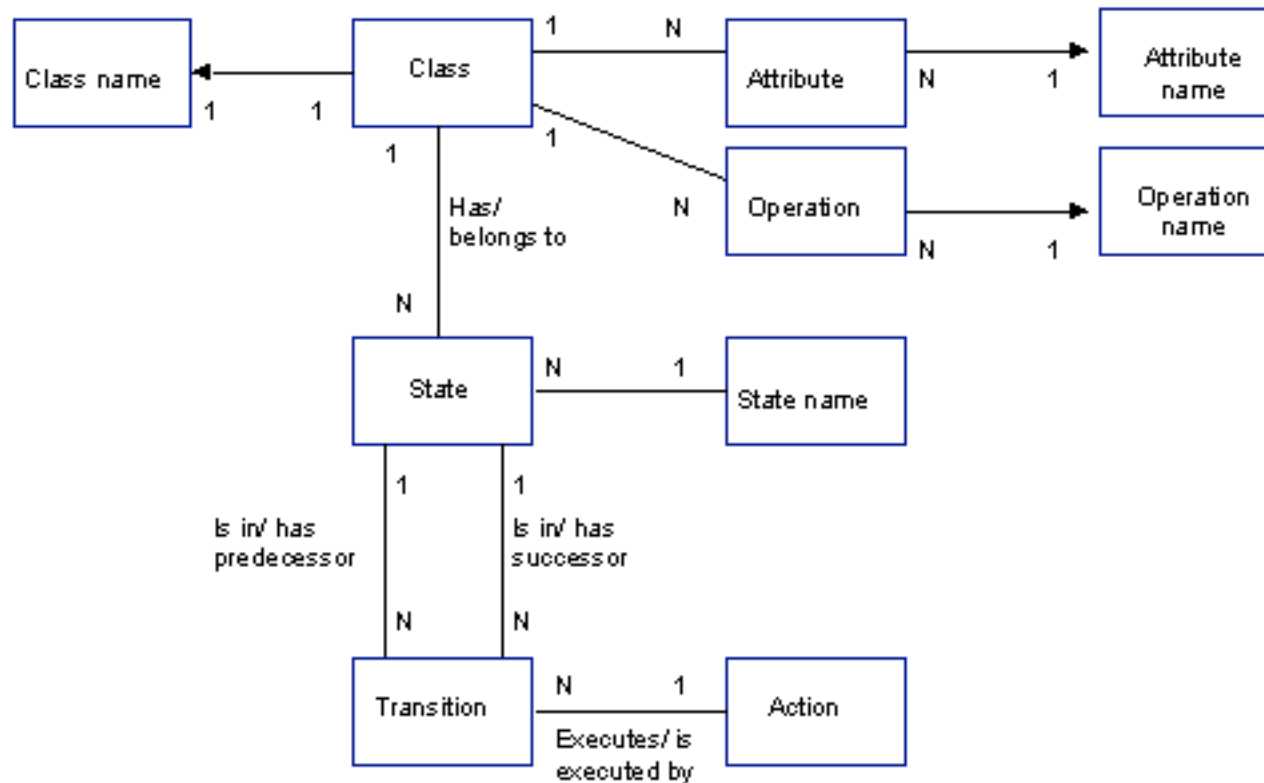
Relational Multiplicities

Examples of relational multiplicities in UML...



Relational Multiplicities

Examples of relational multiplicities in Entity-Relation models



Relational Multiplicities in Alloy

Multiplicities in Alloy can be placed on the elements of a relation declaration...

$$r : S \ m \rightarrow \ n \ T$$

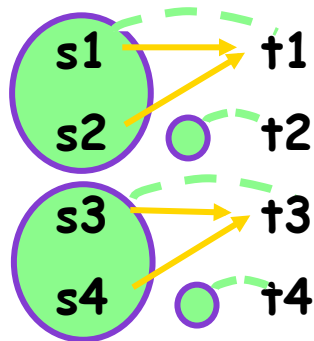
- A multiplicity m on the left (domain) tells you the size of the set associated with each range set element.
- A multiplicity n on the right (range) tells you the size of the set associated with each domain set element.

2 from S
for t1

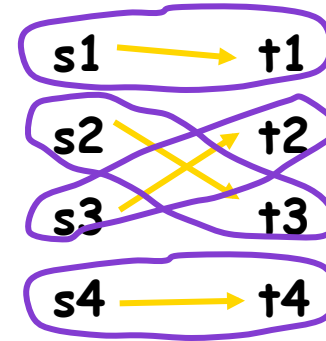
0 from S
for t2

2 from S
for t3

0 from S
for t4



Multiplicity: set S



1 from T
for s1

1 from T
for s3

1 from T
for s2

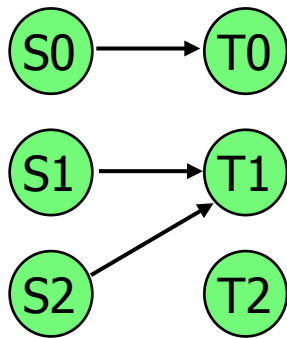
1 from T
for s4

Multiplicity: one T

Functions and Injections

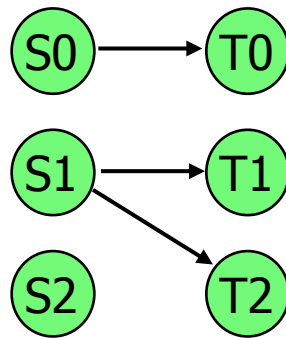
Examples

S \rightarrow one T



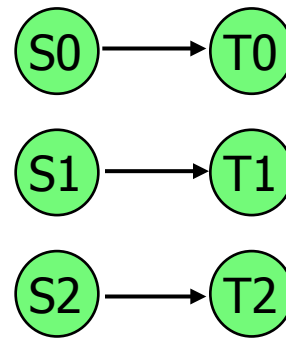
*functional,
but not injective*

S one \rightarrow T



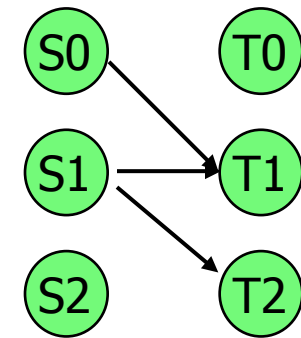
*injective,
but not functional*

S one \rightarrow one T



*functional,
and injective*

S \rightarrow T



*neither functional,
nor injective*

Note: due to the "at most one" in the definitions below, it would be valid to replace the instances of **one** above with **lone**

- A binary relation that maps each atom to at most one other atom is said to be *functional*, and is called a *function*.
- A binary relation that maps at most one atom to each atom is *injective*.

Common Relations via Multiplicities

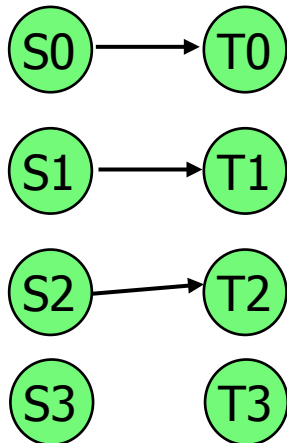
Examples

- $r: A \rightarrow \text{one } B$
 - A (total) function with domain A , range B
- $r: A \text{ one} \rightarrow B$
 - An injective relation
- $r: A \rightarrow \text{lone } B$
 - A (partial) function
- $r: A \text{ one} \rightarrow \text{one } B$
 - An injective (total) function
- $r: A \text{ some} \rightarrow \text{some } B$
 - A surjective relation

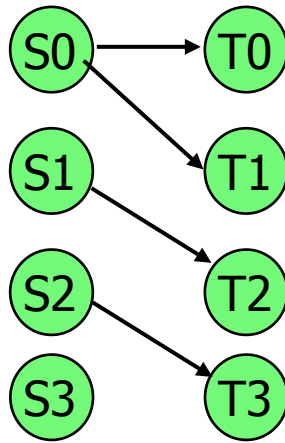
For You To Do

- For each of the relations, give the most precise declaration with multiplicities.

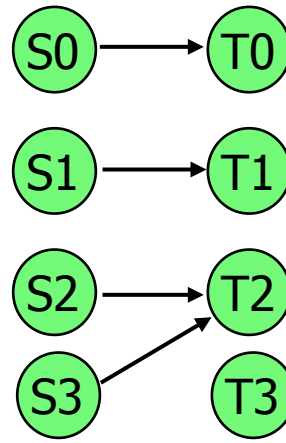
(A)



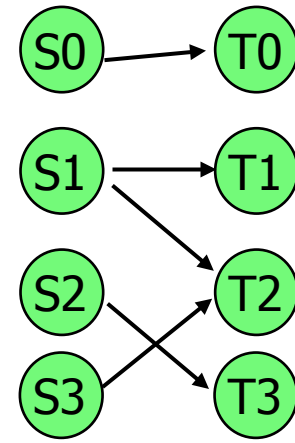
(B)



(C)



(D)



- For each of the declarations below, draw two relations with the same domain / range as the relations above such that the first relation satisfies the multiplicities in the declaration while the second one violates it.
 - r: S some -> lone T
 - r: S set -> one T
 - r: S lone -> lone T
 - r: S some -> some T

Multiplicities are a Shorthand

Multiplicities are just a shorthand, and can be replaced by standard constraints...

$r : A \ m \rightarrow n \ B$

can be written as...

all $a : A \mid n \ a.r$

all $b : B \mid m \ r.b$

Example

$members : Group \ \mathbf{lone} \rightarrow \ \mathbf{some} \ Addr$

can be written as...

all $g : Group \mid \ \mathbf{some} \ g.members$

all $a : Addr \mid \ \mathbf{lone} \ members.a$

Generalizing to Tuples

In the schema below, A and B can be arbitrary expression, and don't have to be relation names.

$r : A \ m \rightarrow \ n \ B$

...in such a case, this says that r maps m tuples in A to each tuple in B , and maps each tuple in A to n tuples in B .

Example

$addr : (\text{Book} \rightarrow \text{Name}) \rightarrow \mathbf{1one} \ \text{Addr}$

...says that that relation $addr$ associates at most one address with each address book / name pair.

✓ { (B0, N0, A1)
(B0, N1, A2)
(B1, N0, A3) }

✗ { (B0, N0, A1)
(B0, N0, A2)
(B1, N0, A3) }

Declaration Constraints

Declaration syntax can also be used to impose constraints on relations that have already be declared, or on arbitrary expressions...

Original declaration...

```
address: (Group + Alias) -> Addr
```

...imposing additional constraints somewhere later in the model...

```
(Alias <: address): Alias -> lone Addr
```

...says that each alias maps to at most one address

Declaration constraints, like any other formula, can be combined with logical operators, placed inside the body of quantifications, etc.

```
all b: Book | b.addr: Name lone -> Addr
```

...says that each address book is injective (maps at most one name to an address)

Nested Multiplicities

Multiplicities can be nested

A declaration of the form...

`r: A -> (B m -> n C)`

...means that for each tuple in A, the corresponding tuples in B -> C form a relation with the given multiplicity. In the case that A is a set, the multiplicity constraint is equivalent to...

`all a: A | a.r : B m -> n C`

Example

`addr: Book -> (Name lone -> Addr)`

...says that, for any book, each address is associated with at most one name, and is equivalent to...

`all b: Book | b.addr: Name lone -> Addr`

...whereas...

`addr: (Book -> Name) lone -> Addr`

...says that each address is associated with at most book/name combination. The first (but not the second) allows an address to appear in more than one book

For You To Do

- For each of the relation declarations below, give two relation instances such that the first relation satisfies the declaration while the second relation violates the declaration...
 - `addr: (Book -> Name) -> some Addr`
 - `addr: lone Book -> (Name -> Addr)`
 - `addr: one Book -> (one Name -> some Addr)`
- For each of the relation declarations above, write a declaration that does not include multiplicities but instead is accompanied by explicit constraint that achieve the same effect as the declarations above
- Given the declaration below, write a declaration constraint that enforces the additional property that each group maps to one or more addresses
 - `address: (Group + Alias) -> Addr`

Cardinality Constraints

The operator `#` applied to a relation gives the number of tuples it contains, as an integer value.

```
Address: (Group + Alias) -> Addr
```

```
...
```

```
all g: Group | #g.address > 1
```

...says that every group has more than one address associated with it

```
addr: Book -> Name -> Addr
```

```
...
```

```
all a: Addr | #addr.a <= 5
```

...says that the number of book/name pairs associated with each address is less than or equal to 5

Sum Expression

The expression below denotes the integer obtained by summing the values of the integer expression ie for all values of the scalar drawn from the set e .

```
sum x: e | ie
```

Example

```
addr: Book -> Name -> Addr
```

```
sum b : { b:Book | #b.addr > 5 } | #b.addr
```

...gives the total number of entries across all books that have great than five entries

Acknowledgements

- The material in this lecture is based on Sections 3.6 and 3.7 from...
 - *Software Abstractions: Logic, Language, and Analysis*, Daniel Jackson, MIT Press, 2006.