

# CIS 771: Software Specifications

## Lecture 14: Advanced OCL Expressions

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## Outline

- Coding transitive closure with recursion
- Useful expressions
- Undefined values
- Meta-modeling

*...with the Academia model as the running example.*

## Transitive Closure in OCL

- OCL does not have a primitive operation for transitive closure
- OCL does allow recursion
- We must implement transitive closure directly in terms of recursion

## Transitive Closure in OCL

Consider the following definitions (transitive-closure-1.use)

```
class A
end
association R between
  A role pred
  A role succ
end
```

We can attempt to code the transitive closure of R as follows

```
class A
operations
  closure() : Set(A) =
    succ.closure()->asSet()->including(self)
end
```

# Transitive Closure in OCL

Consider the following instantiation  
(transitive-closure-instantiation-1.cmd)

```
!create a1:A
!create a2:A
!create a3:A
!insert (a1, a2) into R
!insert (a2, a3) into R
```

An example evaluation

```
use> ? a1.closure()
-> Set{@a1, @a2, @a3} : Set(A)
```

# Transitive Closure in OCL

What is happening on a1.closure?

```
class A
operations
  closure() : Set(A) =
    succ.closure() ->asSet() ->including(self)
end
```

Tracing the evaluation through the recursion...

```
Level 1 call: self = a1, a1.succ = a2
Level 2 call: self = a2, a2.succ = a3
Level 3 call: self = a3, a3.succ = {}
Level 3 return: Set{@a3}
Level 2 return: Set{@a2, @a3}
Level 1 return: Set{@a1, @a2, @a3}
```

## For You To Do...

- Pause the lecture...
- Load the model in *transitive-closure-1.use* into USE
- Run the script  
*transitive-closure-instantiation-1.cmd*
- Now give the following command at the USE command line  
`use> ? a1. closure()`
  - what happens?
- Now give the following commands at the USE command line  
`use> !insert (a3, a1) into R`  
`use> ? a1. closure()`
  - what happens? why? can you fix the problem?

## Transitive Closure in OCL

Consider the following instantiation  
(*transitive-closure-instantiation-2.cmd*)

---

```
!create a1: A
!create a2: A
!create a3: A
!insert (a1, a2) into R
!insert (a2, a3) into R
!insert (a3, a1) into R
```

An example evaluation

---

```
use> ? a1. closure()
..java.lang.RuntimeException: StackOverflow...
```

## Assessment

- The problem is that we have an infinite path through  $R$  and the *closure* operation doesn't know how to stop.
- Intuitively, we should stop when we have collected all the elements that we encounter when walking across  $R$  starting from the initial value (e.g.,  $a1$ ).
- In other words, we should stop when we don't find anything "new" when walking across  $R$ .

## If-then-else

- if *bool-expr* then *expr1* else *expr2* endif
  - Returns *expr1* if *bool-expr* is true
  - Returns *expr2* if *bool-expr* is false
  - Undefined if *bool-expr* is undefined

*...we can use the if-then-else construct to help us code an appropriate transitive closure operation*

# Transitive Closure in OCL

## The correct coding of (reflexive) transitive closure

```
closure(s : Set(A)) : Set(A) =  
  if s->includesAll(s.succ->asSet) then s  
  else closure(s->union(s.succ->asSet))  
  end if
```

Note: stop when we don't find anything new via R (succ) to add to s.

Note: the closure is reflexive because argument s must be included in the result

## An initial call to compute reflexive transitive closure of {self}

```
reachableFromSelf() : Set(A) = closure(Set{self})
```

# Transitive Closure in OCL

## What is happening on `a1.reachableFromSelf()` ?

```
class A  
operations  
  closure(s : Set(A)) : Set(A) =  
    if s->includesAll(s.succ->asSet) then s  
    else closure(s->union(s.succ->asSet))  
    end if  
  reachableFromSelf() : Set(A) = closure(Set{self})  
end
```

## Tracing the evaluation through the recursion...

```
Level 1 call: s = {@a1}, s.succ = {@a2}  
Level 2 call: s = {@a1,@a2}, s.succ = {@a2,@a3}  
Level 3 call: s = {@a1,@a2,@a3}, s.succ = {@a1,@a2,@a3}  
Level 3 return: Set{@a1,@a2,@a3}  
Level 2 return: Set{@a1,@a2,@a3}  
Level 1 return: Set{@a1,@a2,@a3}
```

## For You To Do...

- Pause the lecture...
- Load the model in *transitive-closure-2.use* into USE
- Run the script  
*transitive-closure-instantiation-2.cmd*  
Note that this script adds (a3,a1) to R to create a cycle in R
- Now give the following command at the USE command line  

```
use> ? a1. reachableFromSelf()
```

  - what happens? why?

## Enumeration Types (per OCL spec)

### General Form

enum {value<sub>1</sub>, value<sub>2</sub>, ..., value<sub>n</sub>}

### Example: Academia Grades

enum {A, B, C, D, F, X, W}

### Enumeration Values

#A, #B, #C, #D, #F, #X, #W

## Enumeration Types (per USE)

General Form – declare an *enum* type (e.g., at top of model)

```
enum TypeName {value1, value2, ..., valuen}
```

Example: Academia Grades

```
enum Grade {A, B, C, D, F, X, W}  
...  
class TranscriptEntry  
  attributes  
    course : Course  
    grade  : Grade  
end
```

```
use> create e:TranscriptEntry  
use> !set e.grade = #A
```

## Ordered Associations

- Sometimes we want the result of navigating an association to be a sequence.
- Example:

```
association offspring between  
  Person[0..2] role parents  
  Person[*] role children ordered  
end
```

- Then `p.children` is a sequence.

## Operations on Sequences

- `s->at(i)`            *the  $i$ th element of  $s$*
- `s->first()`            *the first element of  $s$*
- `s->last()`            *the last element of  $s$*
- `s->append(a)`        *adds  $a$  to end*
- `s->prepend(a)`       *adds  $a$  to front*
- `s->asSet()`           *converts to a set*

## let Expressions

- `let  $x$  : Type =  $expr1$  in  $expr2$` 
  - evaluates  $expr2$  with each occurrence of  $x$  replaced by the value of  $expr1$
  - avoids evaluating the same expression multiple times

## Example

```
context Person inv:  
  let income : Integer = self.job.salary->sum in  
  if isUnemployed then  
    income < 100  
  else  
    income >= 100  
  endif
```

## Helper Operations

<pre>... let x : Type1 = expr1 in   ...   ...x...   ...   ...x... ...</pre>	<pre>... f(expr1) ... f(x : Type1) : Type2 =   ...   ...x...   ...   ...x...</pre>
---	--

## For You To Do...

- Pause the lecture...
- Extend the model in *academia-7.use* as follows...
  - This model already contains an extension to *academia-5.use* that adds grades as an enumeration type to a *TranscriptEntry* class as done earlier in the lecture.
  - In the *Transcript* association, declare *transcriptEntries* to be ordered.
  - Using an enumeration type, add a *status* attribute to *Student* that can take on the values *#Normal* or *#Probation*.
  - Write an invariant that says that a student's status is normal iff they only have grades of A's and B's on their transcript. For this invariant, you may want to use a *let* expression since USE has no *iff* construct as a primitive. Specifically, you have to use *implies* twice and reverse the order of the arguments. Use a *let* to avoid duplicating large expressions.
  - Using transitive closure, add an invariant that states that there are no cycles in the prerequisite structure for courses.
  - Write a script to test your extensions.

## Undefined Expressions

- Some expressions that can be undefined
  - `object.oclAsType(T)`
    - ...undefined when type of object has no subtype T
  - `sequence->at(i)`
    - ...undefined when i is greater than length of sequence
  - `sequence->subSequence(i,j)`
    - ...undefined when i,j lie outside the bounds of the sequence or when  $i > j$
  - etc,

## Undefined Expressions

- Undefined expressions tend to propagate
  - if bool-expr then expr-1 else expr-2
    - ...undefined if bool-expr is undefined
  - ...many other examples
- Exceptions:
  - true or anything = true
  - false and anything = false

## For You To Do...

- Pause the lecture...
- Create some expressions whose values are undefined.
- Create some expressions where undefined values are propagated.
- Create some examples where *and* and *or* absorb the undefined values.

## Collections are Flat (per OCL)

- In OCL,  
`Set{Set{1, 2}, Set{2, 3}}`  
and  
`Set {1, 2, 3}`  
have the same value.
- This happens implicitly and is beyond your control.

## Collections Are Usually Not Flat (per USE)

- In USE, Collection types can be nested to any level, e.g.,
  - `Bag(Set(Sequence(Person)))`.
- Implicit flattening is only done when used with the shorthand notation for *collect*.

## Collections Are Usually Not Flat (per USE)

- You can always explicitly flatten a collection with the *flatten* operation that has been added in USE.

- For example,

`company.branches->collect(c | c.employees)`

results in `Bag(Set(Employee))`. This result value can be flattened into a `Bag(Employee)` by using the following expression:

`company.branches->collect(c | c.employees)->flatten`

## For You To Do...

- Pause the lecture...
- Try some examples of nested collections in USE (e.g., you can even use the *transitive-closure* models, and then define collections as literals)
- Flatten them with the *flatten* operation

## Meta Properties

- `type.name` : String
- `type.attributes` : Set(String)
- `type.associationEnds` : Set(String)
- `type.operations` : Set(String)
- `type.supertypes` : Set(OclType)
- `type.allSupertypes` : Set(OclType)
- `type.allInstances` : Set(type)

*Note: it appears that only the last property is supported in USE.*

## Acknowledgements

- Material for this lecture is based on the following sources
  - Chapter 7 (the OCL chapter) of the OMG-UML specification (version 1.3 – March 2000)